

# Mean Field Analysis of Gossip Protocols

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1 Mean Field Analysis

2 An example

# Properties of gossip protocols [BAHKSHI et al. 2007]

## general

- simplicity
- scalability
- symmetry

## functional

- connectivity
- convergence
- shortest path length

## non-functional

- time complexity
- message complexity
- elasticity

## $N \rightarrow \infty$ instead of $N$ large?

- gossip protocols are most suitable if the number of nodes  $N$  is large
- model-based simulation or analytical evaluation of such systems is in general very costly
- is it possible to calculate the limit of the measure of interest for  $N \rightarrow \infty$ ?
- yes — sometimes it is!
- convergence to *mean field*, a deterministic dynamical system which constitutes a limit for  $N \rightarrow \infty$  [Le Boudec et al. 2007]

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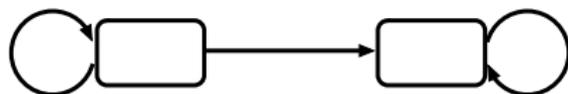
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# Interacting objects

- $N$  interacting objects
- each with discrete (finite) state space  $S$
- time is discrete,  $T = 0, 1, 2, \dots$
- random variable  $X_n^N(t)$  state of object  $n$  at time  $t$



# Occupancy measure & global memory

- random variable  $M_i^N(t)$ : **fraction of objects** that is in state  $i$  at time  $t$

$$M_i^N(t) = \frac{1}{N} \sum_{n=1}^N 1_{\{X_n^N(t)=i\}}$$

- random variable  $\bar{R}^N(t)$ : **history** of the occupancy measure

$$\bar{R}^N(t+1) = g(\bar{R}^N(t), \bar{M}^N(t+1))$$

# Transitions

- transition probabilities for each object depend on
  - ▶ the current state  $i$
  - ▶ the current global memory  $\bar{r}$

$$P_{i,j}^N(\bar{r}) = \Pr\{X_n^N(t+1) = j \mid \bar{R}^N(t) = \bar{r}, X_n^N(t) = i\}$$

$$\lim_{N \rightarrow \infty} P_{i,j}^N(\bar{r}) = P_{i,j}(\bar{r})$$

# Convergence to mean field

assume

- $\overline{M}^N(0)$  converges almost surely to deterministic limit  $\overline{\mu}(0)$ ,
- $\overline{R}^N(0)$  converges almost surely to deterministic limit  $\overline{\rho}(0)$ .

define

$$\left. \begin{aligned} \overline{\mu}(t+1) &= \overline{\mu}(t) \cdot \mathbf{P}(\overline{\rho}(t)), \\ \overline{\rho}(t+1) &= g(\overline{\rho}(t), \overline{\mu}(t+1)). \end{aligned} \right\} \text{mean field}$$

then for any  $t$ , almost surely,

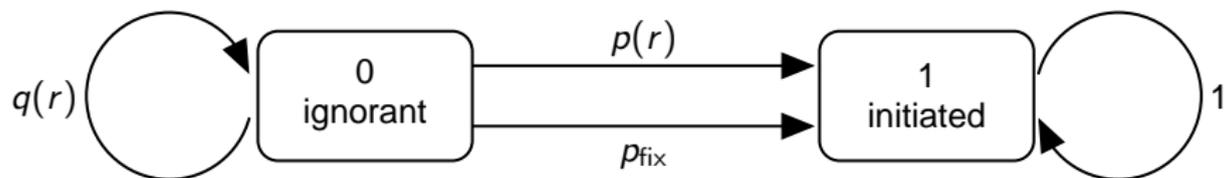
$$\lim_{N \rightarrow \infty} \overline{M}^N(t) = \overline{\mu}(t) \quad \text{and} \quad \lim_{N \rightarrow \infty} \overline{R}^N(t) = \overline{\rho}(t).$$

(almost surely = with probability one)

# An Example

- P2P network with  $N$  nodes
- random overlay, in-degree = out-degree =  $C$
- we want to model the dissemination of a message
- with probability  $p_{\text{fix}}$  a nodes receives the message from outside the network
- if a node is initiated it sends the message to  $K$  nodes in its partial view per time step

## A single node



occupancy measure

$$\bar{M}^N(t) = (M_0^N(t), M_1^N(t))$$

global memory

$$R^N(t) = M_1^N(t+1)$$

transition probabilities

$$\mathbf{P}(r) = \begin{pmatrix} q(r) & p(r) + p_{\text{fix}} \\ 0 & 1 \end{pmatrix}$$

# Transition probabilities

assume:  $N$  nodes,  $r \cdot N$  nodes are infected

- look at a single, non-infected node and take a single in-link
  - ▶ probability that this in-link comes from an ignorant node

$$q_1(r) = \frac{N - rN}{N} = 1 - r$$

- ▶ probability that this in-link comes from an initiated node but no message is sent (since only  $K$  are selected)

$$q_2(r) = \frac{rN}{N} \cdot \frac{C - K}{C} = r \cdot \frac{C - K}{C}$$

- probability that no message reaches this node from another node

$$q'(r) = \sum_{i=0}^C q_1^i \cdot q_2^{C-i}$$

- probability to get infected

$$q(r) = q'(r) - p_{\text{fix}}$$

# Mean field

all nodes are ignorant at the beginning

$$\bar{\mu}(0) = (1, 0)$$

global memory corresponds to fraction of initiated nodes

$$\rho(0) = 0$$

equations for the mean field

$$\bar{\mu}(t+1) = \bar{\mu}(t) \cdot \begin{pmatrix} q(\rho(t)) & 1 - q(\rho(t)) \\ 0 & 1 \end{pmatrix}$$
$$\rho(t+1) = \mu_1(t+1)$$

## Simulation for finite $N$

- modelling each object individually would result in  $2^N$  states
- it suffices to keep track of the number of infected nodes ( $N + 1$  states)
- transition probability from  $i$  to  $j \geq i$  for  $r \cdot N$  infected nodes

$$\Pr(i \rightarrow j) = \binom{N-i}{j-i} \cdot q(r)^{N-j} (1 - q(r))^{j-i}$$

# Time complexity

