

Lazy Data Saturation

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Example: Very simple model of Zeroconf protocol

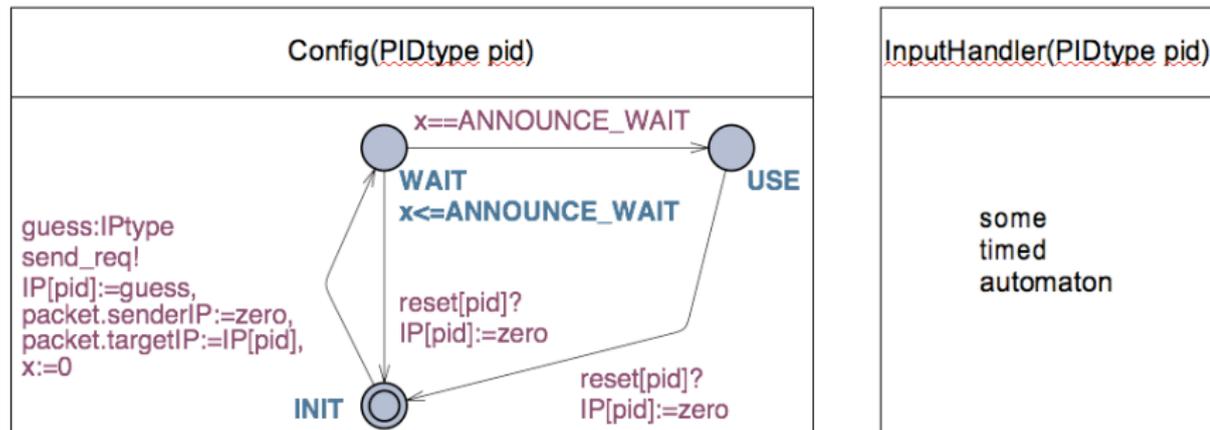
- ▶ Protocol for dynamic configuring IP addresses
- ▶ A large number of hosts communicate via broadcasts
- ▶ Property: *Hosts should not use the same IP address*
- ▶ Main simplifications
 - ▶ No probabilities
 - ▶ No delay in the network

(on my website paper on Zeroconf)

Model of a single host

Globals

```
typedef scalar[k] PIDtype;  
typedef scalar[m+1] Iptype;  
  
typedef struct{  
    Iptype senderIP; // sender IP address  
    Iptype targetIP; // target IP address  
} ARP_packet;  
  
Iptype zero;  
ARP_packet packet;
```



Restrictions on use of scalarsets [ID93]

- ▶ The elements can only be compared on equality, no addition or other operators
- ▶ Every scalarset is its own type incomparable to other scalarset types
- ▶ Iterating over a scalarset variable is order independent

Symmetry Reduction on scalar PIDtype

For explanation: **no** InputHandler automata, and **no** clocks

For example $k = 3$ hosts

State vectors:

$$\begin{array}{l} \text{Config}(h_1).\text{loc} \\ \text{Config}(h_1).\text{guess} \\ \text{Config}(h_2).\text{loc} \\ \text{Config}(h_2).\text{guess} \\ \text{Config}(h_3).\text{loc} \\ \text{Config}(h_3).\text{guess} \\ \text{packet.senderIP} \\ \text{packet.targetIP} \end{array} \left(\begin{array}{c} \text{INIT} \\ \text{zero} \\ \text{WAIT} \\ \text{ip}_1 \\ \text{USE} \\ \text{ip}_2 \\ \text{zero} \\ \text{ip}_1 \end{array} \right) \approx \left(\begin{array}{c} \text{INIT} \\ \text{zero} \\ \text{USE} \\ \text{ip}_2 \\ \text{WAIT} \\ \text{ip}_1 \\ \text{zero} \\ \text{ip}_1 \end{array} \right)$$

Symmetry Reduction on scalar IPtype

State vectors:

$$\begin{array}{l} \text{Config}(h_1).\text{loc} \\ \text{Config}(h_1).\text{guess} \\ \text{Config}(h_2).\text{loc} \\ \text{Config}(h_2).\text{guess} \\ \text{Config}(h_3).\text{loc} \\ \text{Config}(h_3).\text{guess} \\ \text{packet.senderIP} \\ \text{packet.targetIP} \end{array} \left(\begin{array}{c} \text{INIT} \\ \text{zero} \\ \text{WAIT} \\ \text{ip}_1 \\ \text{USE} \\ \text{ip}_2 \\ \text{zero} \\ \text{ip}_1 \end{array} \right) \approx \left(\begin{array}{c} \text{INIT} \\ \text{zero} \\ \text{WAIT} \\ \text{ip}_2 \\ \text{USE} \\ \text{ip}_1 \\ \text{zero} \\ \text{ip}_2 \end{array} \right)$$

Data saturation [ID93]

Suppose the number of hosts is fixed to 3.

- ▶ Verification makes no difference for more than 6 IP addresses available.

Intuition: Simply look at the state vector.

- ▶ **But** behaviour shows that at most 3 IP addresses are needed.

Intuition: packet.senderIP and packet.targetIP are only set to IPs already in use.

IPtype is a so called **data scalar**.

Data scalarsets

Extra conditions:

- ▶ not used as array indices
- ▶ not used for template arguments
- ▶ not used in for-statements

Check all instances?

In the example: data saturation for > 6 #IPs

All instances of the model (#IP = 1, = 2, ..., = 6) must be verified.

Example



Idea: use only guards having only existential quantification in their Negation Normal Form

Lazy Data Saturation

In *selection* and *quantification* over data scalars, elements of the scalar are chosen.

Idea: either

1. pick a value that is not used anywhere in the state vector
2. or pick some used value

State vectors:

$$\begin{array}{l} \text{Config}(h_1).\text{loc} \\ \text{Config}(h_1).\text{guess} \\ \text{Config}(h_2).\text{loc} \\ \text{Config}(h_2).\text{guess} \\ \text{Config}(h_3).\text{loc} \\ \text{Config}(h_3).\text{guess} \\ \text{packet.senderIP} \\ \text{packet.targetIP} \end{array} \quad \left(\begin{array}{c} \text{INIT} \\ \text{zero} \\ \text{WAIT} \\ \text{ip}_1 \\ \text{USE} \\ \text{ip}_2 \\ \text{zero} \\ \text{ip}_2 \end{array} \right) \rightarrow \left(\begin{array}{c} \text{WAIT} \\ ? \\ \text{WAIT} \\ \text{ip}_1 \\ \text{USE} \\ \text{ip}_2 \\ \text{zero} \\ ? \end{array} \right)$$

Future Work

- ▶ Implement for Uppaal model checker
- ▶ Investigate if relaxation of *Diagonal property* is possible.



C.N. Ip and D.L. Dill.

Better verification through symmetry.

In David Agnew, Luc J. M. Claesen, and Raul Camposano, editors, *Proceedings of the 11th IFIP WG10.2 International Conference on Computer Hardware Description Languages and their Applications - CHDL '93*, Ottawa, Ontario, Canada, 26-28 April, 1993, volume A-32 of *IFIP Transactions*, pages 97–111. North-Holland, 1993.